#### **KVM and Hypervisor Security**

David Shepard and Matt Gaston CMU/SEI Cyber Innovation Center

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### **Overview**

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Dynamic On-Demand High-Performance Computing System

KVM and Hypervisor Security

Recommendations

Acknowledgements

### **CMU/SEI Cyber Innovation Center**

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- Shape and leverage academic and industrial research wherever possible
- Employ creative solutions to finding, assessing, and proving technology capabilities for mission applications
- Develop and extend software technologies tailoring them for application to government mission needs
- Promote government awareness and knowledge of emerging technologies and their applications



# **Motivation: A Heterogeneous HPC Utility Cloud**

#### Dynamic On-Demand High-Performance Computing System (DODCS)





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### **Kernel-based Virtual Machine Architecture**



Source: Linux Information - Virtualization on Linux - KVM, IBM, 2011

# BLUF

It is possible to secure a KVM-based hypervisor against an appropriate risk target for a defined level of security.

Compelling advantages of KVM:

- The majority of the footprint is running as an unprivileged user-mode process
- Draws on a huge base of drive support from the Linux kernel
- Open source

When considering hypervisor security, the Type 1 vs. Type 2 distinction is not helpful. Rather, the focus should be on security aspects of virtualization technology and how specific implementations address these aspects.

Note: There is no such thing as perfect security.



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**Performance:** KVM uses hardware support for virtualization, effectively running on 'bare metal.' KVM allows the guest VM to run at processor ring level zero for performance.

#### Security:

- Type 1 hypervisors make use of privileged guest VMs for maintenance and management – a compromise of a guest compromises the hypervisor
- KVM uses privilege "de-escalation" there is no privileged guest
  VM only user requests that require privilege escalation use it
- The KVM software stack is minimal only the kernel, a few system daemons, and QEMU (comparable to Xen)

### **The Threat Model**

The threat considered: "regular" users – insiders

Types of security breaches:

- VM Escape compromise of the hypervisor and assumption of control over all VMs
- **Privilege Escalation** an exploit that allows an unprivileged VM to execute code in a privileged context
- **Denial of Service** system crash or access to system resources are denied

# **Known Vulnerabilities**

Vulnerability databases:

- NIST National Vulnerability Database
- MITRE Common Vulnerabilities and Exposures
- USCERT Vulnerability Database



# of recent vulnerabilities

As of June 1, 2011



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# An Example KVM Vulnerability



### **General Recommendations**

- Apply a proper SELinux security policy to strengthen resource separation.
- Patch, patch, patch.
- Use physical separation for different authorities.
- Keep logs of guest behavior
- Keep logs of administrative activities
- Backup critical systems.

# **Security Considerations for KVM**

- It does not appear unsafe to support access to physical hardware in support of provisioning heterogeneous processors as the access is limited to the VM guest.
- More robust to failure since guest VMs run as unprivileged processes
- To compromise the "host" from a guest VM (unprivileged) is exceedingly difficult
- Small code footprint

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And Remember, there is no such thing as perfect security.



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